

## Lecture 2: Separation and Duality

TTIC 31070 / CMSC 35470 / BUSF 36903 / STAT 31015

Convex Optimization

Prof. Zhiyuan Li

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## The Plan for Lecture 2

Lecture 1 said convexity makes local information global.

Lecture 2 turns that slogan into one geometric picture:

A closed convex set is exactly the intersection of all halfspaces that contain it.

The route:

1. Encode convex functions as sets via the **epigraph**
2. Develop **projection** and **separation** for convex sets
3. Translate back to functions through **subgradients**

## Extended Real Line and Effective Domain

**Definition 2.1** (Extended real line). We write  $\mathbb{R} \cup \{+\infty\}$  for the one-sided extended real line, with  $a \leq +\infty$  for every  $a$ , and the conventions

$$a + (+\infty) = +\infty, \quad \lambda(+\infty) = +\infty \ (\lambda > 0), \quad 0 \cdot (+\infty) := 0, \quad \inf \emptyset = +\infty.$$

**Definition 2.2** (Effective domain). For  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$ , the effective domain is

$$\text{dom } f := \{x \in E : f(x) < +\infty\}.$$

The value  $+\infty$  encodes “outside the domain.”

## Convex Extended-Value Function

**Definition 2.3** (Convex extended-value function).  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  is **convex** if

$$\forall x, y \in E, \forall \theta \in [0, 1], \quad f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y).$$

This is the same inequality as in Lecture 1, but now the domain is encoded inside  $f$  via  $+\infty$ : if  $h : C \rightarrow \mathbb{R}$  is convex, define  $\tilde{h} = h$  on  $C$  and  $\tilde{h} = +\infty$  outside  $C$ . Then  $h$  is convex iff  $\tilde{h}$  is.

Lecture 2 does not change the notion of convexity. It only changes the packaging.

# Epigraph

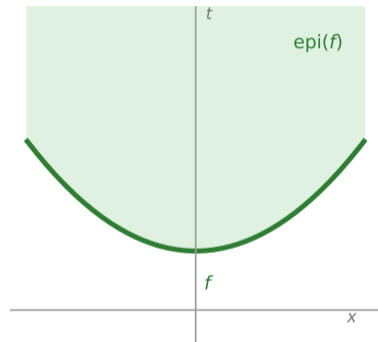
**Definition 2.4** (Epigraph). For  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$ :

$$\text{epi}(f) := \{(x, t) \in E \times \mathbb{R} : t \geq f(x)\}.$$

**Lemma 2.1** (Epigraph criterion).  $f$  is convex  $\iff$   $\text{epi}(f)$  is a convex set.

**Proof idea.** Convex combination of two points above the graph stays above the graph.  $\square$

After this, we leave functions and work with convex **sets** directly.



# From Functions to Sets

After the epigraph criterion, the lecture switches to convex **sets**.

The first geometric input is **Euclidean projection**. Projection is what manufactures separating hyperplanes, and later those separating hyperplanes will come back as subgradients.

The chain:

projection  $\longrightarrow$  separation  $\longrightarrow$  support  $\longrightarrow$  subgradients

# Halfspace, Interior, and Closure

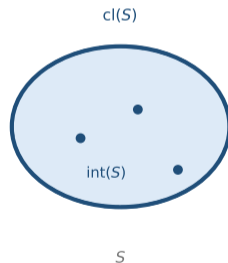
**Definition 2.5** (Closed halfspace). For  $\xi \in E^* \setminus \{0\}$  and  $b \in \mathbb{R}$ :

$$H(\xi, b) := \{x \in E : \langle \xi, x \rangle \leq b\}.$$

**Definition 2.6** (Interior and closure).

$$\text{int}(S) := \{x \in S : \exists r > 0, B(x, r) \subseteq S\}$$

$$\text{cl}(S) := \bigcap \{F : F \text{ closed}, S \subseteq F\}$$



$\text{int}(S) \subseteq S \subseteq \text{cl}(S)$ .  $S$  open  $\Leftrightarrow \text{int}(S) = S$ ;  $S$  closed  $\Leftrightarrow \text{cl}(S) = S$ .

## Projection onto a Closed Convex Set

**Theorem 2.3** (Projection). Let  $C \subseteq E$  be nonempty, closed, and convex, with auxiliary Euclidean norm  $\|\cdot\|_2$ . For every  $z \in E$  there exists a unique  $p \in C$  such that

$$\|z - p\|_2 = \inf_{y \in C} \|z - y\|_2.$$

Moreover:  $p$  is the projection iff  $\langle z - p, y - p \rangle_2 \leq 0$  for all  $y \in C$ .

**Proof idea.** Existence: Weierstrass on  $C \cap \overline{B}(z, R)$ . Uniqueness: strict convexity of  $\|\cdot\|_2^2$ . Characterization: first-order condition for the constrained minimum.  $\square$

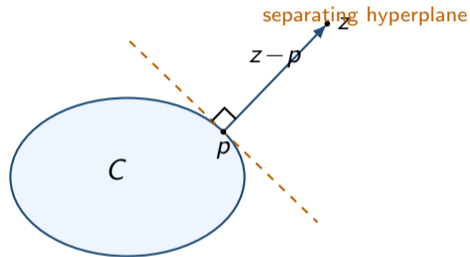
# Point-to-Set Separation

How do we actually separate a point from a closed convex set? Project, and use the normal vector.

**Theorem 2.2** (Point-to-set separation). Let  $C \subseteq E$  be nonempty, closed, and convex,  $z \notin C$ . Then  $\exists \xi \in E^* \setminus \{0\}, b \in \mathbb{R}$ :

$$\langle \xi, z \rangle > b \quad \text{and} \quad \forall y \in C, \langle \xi, y \rangle \leq b.$$

**Proof idea.** Project  $z$  onto  $C$ ; set  $\xi := z - p$ ,  $b := \langle \xi, p \rangle$ .  $\square$



## Why Is Closedness Needed?

**Question.** Can we drop the closedness assumption in the separation theorem?

**No.** Let  $C = (0, 1) \subseteq \mathbb{R}$  and  $z = 1$ .

Then  $z \notin C$ , but no strict separating halfspace exists: any halfspace  $\{x : x \leq b\}$  that contains  $C$  must have  $b \geq 1$ , so it also contains  $z$ .

Closedness is needed both for the **proof** (projection may not exist) and for the **statement** (separation can genuinely fail).

## Separation as a Certificate

The separation theorem gives a **certificate** that  $z \notin C$ :

- ▶ To *prove*  $z \in C$ , one must exhibit a representation of  $z$  as a feasible point.
- ▶ To *prove*  $z \notin C$ , it suffices to exhibit a halfspace  $H(\xi, b)$  with  $C \subseteq H(\xi, b)$  and  $z \notin H(\xi, b)$ .

This is the **certificate viewpoint** that recurs throughout the course: duality provides short, checkable proofs of infeasibility and optimality.

Lecture 3: Farkas' lemma as an infeasibility certificate for LP.

Lecture 5: KKT multipliers as an optimality certificate.

## Example: The PSD Cone as a Proof System

**Example.** Work in symmetric matrices  $\mathbb{S}^n$  with Frobenius pairing  $\langle Y, Z \rangle := \text{tr}(YZ)$ . Then:

$$\mathbb{S}_+^n = \bigcap_{u \in \mathbb{R}^n} \{X \in \mathbb{S}^n : \langle uu^\top, X \rangle \geq 0\}.$$

If  $X \notin \mathbb{S}_+^n$ , there exists  $u$  with  $u^\top Xu < 0$ .

So  $uu^\top$  is a **concrete certificate** that  $X$  is not PSD.

The separation theorem says this kind of linear proof system is **complete**: every false claim “ $X \succeq 0$ ” can be refuted by a separating inequality.

## Supporting Hyperplane Theorem

Projection gives separation for **exterior** points. What about **boundary** points? Let  $z_k \rightarrow x$  approach the boundary and pass to the limit.

**Theorem 2.4** (Supporting hyperplane). Let  $C \subseteq E$  be nonempty and convex,  $x \in \text{cl}(C) \setminus \text{int}(C)$ . Then there exist  $\xi \in E^* \setminus \{0\}$  and  $b \in \mathbb{R}$  such that

$$\langle \xi, x \rangle = b \quad \text{and} \quad \forall y \in C, \quad \langle \xi, y \rangle \leq b.$$

**Proof idea.** Take a sequence  $z_k \rightarrow x$  with  $z_k \notin C$ . Separate each  $z_k$  from  $C$ , normalize the separating covectors on the unit sphere, extract a convergent subsequence (compactness), and pass to the limit.  $\square$

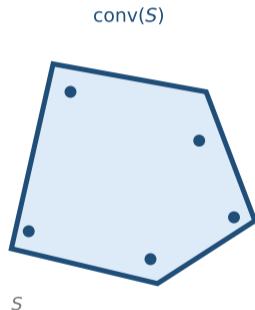
# Convex Hull and Closed Convex Hull

**Definition 2.7** (Convex hull).

$$\text{conv}(S) := \bigcap \{C \subseteq E : C \text{ convex}, S \subseteq C\}.$$

**Definition 2.8** (Closed convex hull).

$$\text{cl conv}(S) := \bigcap \{C \subseteq E : C \text{ closed, convex}, S \subseteq C\}.$$



## Basic Properties of Closure, Convex Hull, and Closed Convex Hull

**Proposition 2.5.** Let  $S \subseteq E$ . Then:

1.  $\text{cl}(S)$  is closed,  $\text{conv}(S)$  is convex,  $S \subseteq \text{cl}(S) \cap \text{conv}(S)$ .
2.  $\text{cl conv}(S) = \text{cl}(\text{conv}(S))$ .
3.  $\text{cl conv}(S)$  is closed and convex, and  $S \subseteq \text{cl conv}(S)$ .
4. If  $C$  is closed and convex, then  $\text{cl conv}(C) = C$ .
5.  $\text{cl conv}(\text{cl conv}(S)) = \text{cl conv}(S)$ .

The proof is left as an exercise.

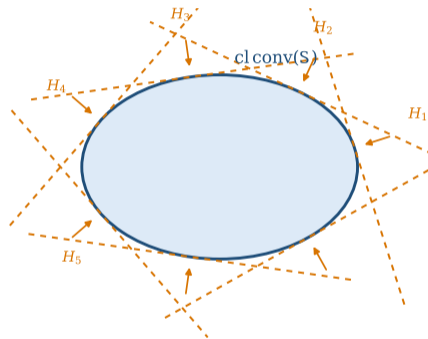
# Closed Convex Hull = Intersection of Halfspaces

Once supporting hyperplanes are available, can we reconstruct a closed convex set from **all** the halfspaces that contain it?

**Theorem 2.6.** For any  $S \subseteq E$ :

$$\text{cl conv}(S) = \bigcap \{H(\xi, b) : S \subseteq H(\xi, b)\}.$$

**Proof idea.**  $\subseteq$ : trivial.  $\supseteq$ : if  $z \notin \text{cl conv}(S)$ , separate.  $\square$



Intersecting all containing halfspaces recovers  $\text{cl conv}(S)$ .

## Convex Cones and Duality

The halfspace-hull picture works for general sets. Cones are the special case where the bias term  $b$  vanishes, so halfspaces become **half-spaces through the origin**.

**Definition 2.9** (Convex cone, dual cone, polar cone).  $K \subseteq E$  is a convex cone if

$$\forall x_1, x_2 \in K, \forall \alpha_1, \alpha_2 \geq 0, \quad \alpha_1 x_1 + \alpha_2 x_2 \in K.$$

Its dual and polar cones are

$$K^* := \{\xi \in E^* : \forall x \in K, \langle \xi, x \rangle \geq 0\}, \quad K^\circ := \{\xi \in E^* : \forall x \in K, \langle \xi, x \rangle \leq 0\}.$$

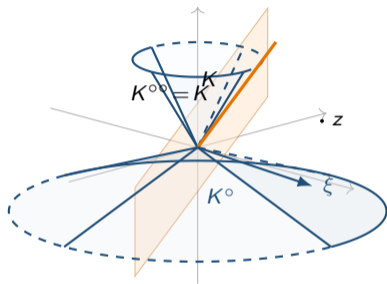
# Cone Separation and Bipolarity

**Theorem 2.7** (Cone separation and bipolarity). Let  $K$  be a nonempty closed convex cone,  $z \notin K$ . Then  $\exists \xi \in K^\circ: \langle \xi, z \rangle > 0$ .  
Consequently,  $K^{\circ\circ} = K$ .

**Proof idea.**

**Separation:** Apply point-to-set separation to  $K$ . Since  $0 \in K$ , we get  $b \geq 0$ . Since  $K$  is a cone,  $\langle \xi, x \rangle \leq 0$  for all  $x \in K$  (otherwise scale  $x$  to violate  $b$ ). So  $\xi \in K^\circ$ .

**Bipolarity:**  $K \subseteq K^{\circ\circ}$  is immediate. If  $z \notin K$ , separation gives  $\xi \in K^\circ$  with  $\langle \xi, z \rangle > 0$ , so  $z \notin K^{\circ\circ}$ .  $\square$



## Affine Sets, Affine Hull, and Relative Interior

The cone theorem is the last set-level result. Before returning to functions, we need **relative interior**: the subgradient existence theorem lives on  $\text{ri}(\text{dom } f)$ , not  $\text{dom } f$ .

**Definition 2.10** (Affine set).  $A \subseteq E$  is **affine** if  $\forall x, y \in A, \forall \theta \in \mathbb{R}, \theta x + (1 - \theta)y \in A$ .

**Definition 2.11** (Affine hull and relative interior).

$$\text{aff}(S) := \bigcap \{A \subseteq E : A \text{ is affine and } S \subseteq A\}.$$

If  $C \subseteq E$ , its relative interior is

$$\text{ri}(C) := \{x \in C : \exists r > 0, \{y \in \text{aff}(C) : \|y - x\| < r\} \subseteq C\}.$$

**Why affine?** Convexity and separation are fundamentally affine notions — no preferred origin is needed. We work in  $E$  but pass locally to affine subspaces when needed.

## Subgradient

We now translate the geometry back into function language. A supporting hyperplane to  $\text{epi}(f)$  through  $(x, f(x))$  is exactly a global affine lower bound on  $f$  — and that is what a subgradient is.

**Definition 2.12** (Subgradient and subdifferential). Let  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$ ,  $x \in \text{dom } f$ . A covector  $g \in E^*$  is a **subgradient** of  $f$  at  $x$  if

$$\forall y \in E, \quad f(y) \geq f(x) + \langle g, y - x \rangle.$$

The set of all such covectors is  $\partial f(x)$ .

**Lemma 2.8** (Subgradients = supporting hyperplanes of  $\text{epi}$ ).

$$g \in \partial f(x) \iff (g, -1) \text{ supports } \text{epi}(f) \text{ at } (x, f(x)).$$

**Proof idea.** Unpack both definitions; they say the same thing.  $\square$

## Proper Function

**Definition 2.13** (Proper extended-value function).  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  is **proper** if

$$\text{dom } f \neq \emptyset.$$

Equivalently, there exists  $x \in E$  with  $f(x) < +\infty$ .

Properness rules out the trivial case  $f \equiv +\infty$ . The subgradient existence theorem requires properness.

## Differentiable Convex $\Rightarrow$ Unique Subgradient

**Proposition 2.9.** Let  $f : E \rightarrow \mathbb{R}$  be convex and differentiable. Then for every  $x \in E$ ,

$$\partial f(x) = \{\nabla f(x)\}.$$

So for smooth convex functions, the subgradient is just the gradient. Subgradients become genuinely new only for **nonsmooth** convex functions.

## Existence of Subgradients

**Theorem 2.10** (Subgradient existence on the relative interior). Let  $f : E \rightarrow \mathbb{R} \cup \{+\infty\}$  be proper and convex. Then

$$\forall x \in \text{ri}(\text{dom } f), \quad \partial f(x) \neq \emptyset.$$

**Proof idea.** The point  $(x, f(x))$  is a boundary point of  $\text{epi}(f)$  inside  $\text{aff}(\text{dom } f) \times \mathbb{R}$ . Apply supporting hyperplane there and read the result as a subgradient.  $\square$

Can we weaken  $x \in \text{ri}(\text{dom } f)$  to  $x \in \text{dom } f$ ?

No:  $f(x) = x \log x$  on  $[0, \infty)$  has  $\partial f(0) = \emptyset$ , since  $\log y \geq g$  for all  $y > 0$  is impossible.

# Summary & What's Next

## Today:

- ▶ Epigraph converts function convexity to set convexity
- ▶ Projection  $\rightarrow$  point-to-set separation  $\rightarrow$  supporting hyperplanes
- ▶  $\text{cl conv}(S) = \bigcap$  containing halfspaces
- ▶ Cone separation and bipolarity:  $K^{\circ\circ} = K$
- ▶ Subgradient = supporting hyperplane of the epigraph
- ▶ Subgradients exist on  $\text{ri}(\text{dom } f)$

## Next lecture:

- ▶ Linear programming as the first application of duality and separation
- ▶ Farkas' lemma as a certificate alternative
- ▶ Strong duality and complementary slackness